

Factors Impacting Performance of Multithreaded Triangular Solve

VECPAR'10

June 23, 2010

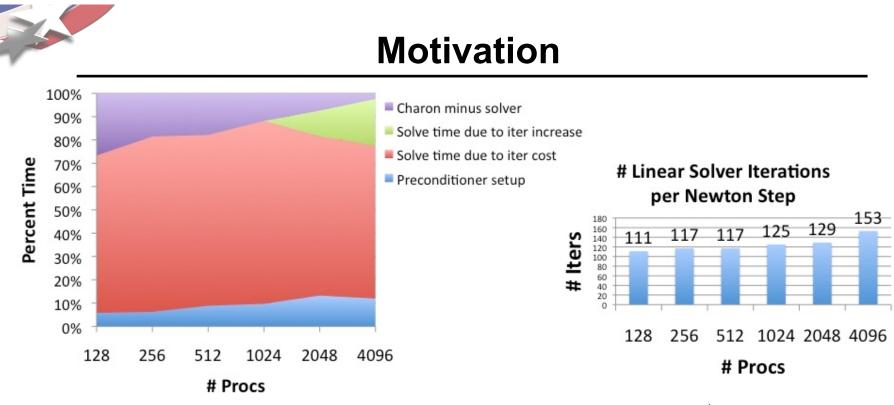
Michael Wolf, Mike Heroux, Erik Boman Extreme-scale Algorithms and Software Institute (EASI)



Sandia National Laboratories is a multi-program laboratory operated by Sandia Corporation, a wholly owned subsidiary of Lockheed Martin company, for the U.S. Department of Energy's National Nuclear Security Administration under contract DE-AC04-94AL85000.



- Triangular solver is important numerical kernel
 - Essential role in preconditioning linear systems
- Difficult algorithm to parallelize
- Trend of increasing numbers of cores per socket
- Threaded or hybrid approach potentially beneficial
- Focus of work: threaded triangular solve on each node/socket

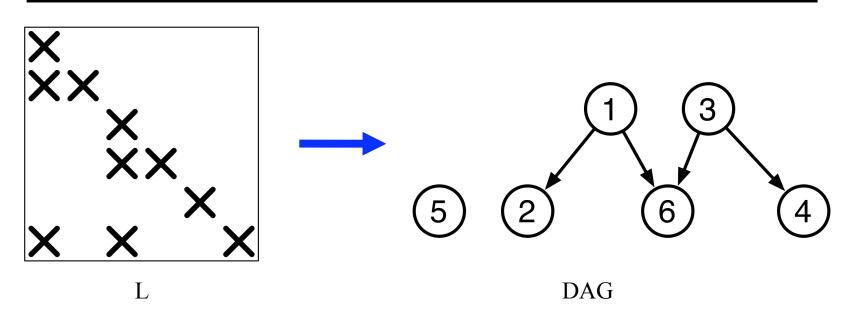


Strong scaling of Charon on TLCC (P. Lin, J. Shadid 2009)

- Inflation in iteration count due to number of subdomains (MPI tasks)
- With scalable threaded triangular solves
 - Solve triangular system on larger subdomains
 - Reduce number of subdomains (MPI tasks)



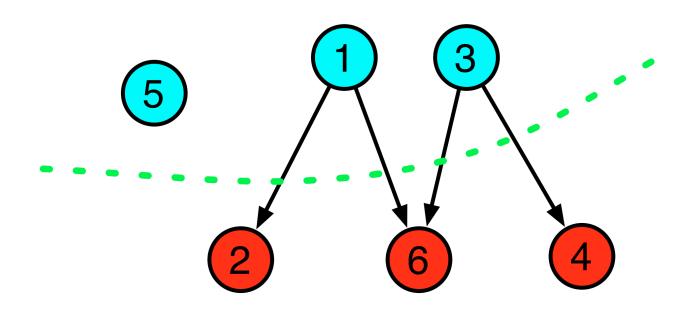
Level Set Triangular Solver



- Initially, focus attention on level set triangular solver (J. Saltz, 1990)
 - Level set approach exposes parallelism
- First, express data dependencies for triangular solve with a directed acyclic graph (DAG)



Level Set Triangular Solver



- Determine level sets of this DAG
 - Represent sets of row operations that can be performed independently



$$\tilde{L} = PLP^{T} = \begin{bmatrix} D_{1} & & & \\ A_{2,1} & D_{2} & & \\ A_{3,1} & A_{3,2} & D_{3} & & \\ \vdots & \vdots & \vdots & \ddots & \\ A_{l,1} & A_{l,2} & A_{l,3} & \dots & D_{l} \end{bmatrix}$$

- Permuting matrix so that rows in a level set are contiguous
 - D_i are diagonal matrices
 - Row operations in each level set can be performed independently



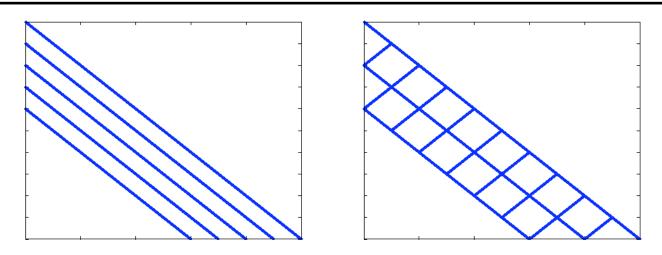
Level Set Triangular Solver

$$\tilde{x}_{1} = D_{1}^{-1} \tilde{y}_{1}
\tilde{x}_{2} = D_{2}^{-1} (\tilde{y}_{2} - A_{2,1} \tilde{x}_{1})
\vdots \vdots \vdots \\
\tilde{x}_{l} = D_{l}^{-1} (\tilde{y}_{l} - A_{l,1} \tilde{x}_{1} - \dots - A_{l,l-1} \tilde{x}_{l-1})$$

- Resulting operations for triangle solve
 - Row operations in each level can be performed independently (parallel for)



Simple Prototype



- Simple prototype of level set threaded triangular solve
 - Assumes fixed number of rows per level
 - Assumes matrices preordered by level
 - Pthreads
- Allowed us to explore factors affecting performance
- Run experiments on two platforms
 - Intel Nehalem: two 2.93 GHz quad-core Intel Xeon processors
 - AMD Istanbul: two 2.6 GHz six-core AMD Opteron processors



Factor 1: Type of Barrier

Algorithm 1 Passive Barrier.

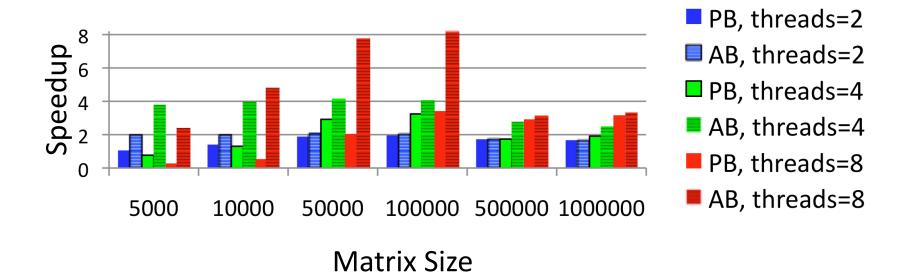
```
void passiveBarrier()
{
    pthread_mutex_lock(&mutex);
    numArrived++;
    if(numArrived < NUM_THREADS) {
        pthread_cond_wait(&barrCond,&mutex);
    }
    else {
        pthread_cond_broadcast(&barrCond);
        numArrived = 0;
    }
    pthread_mutex_unlock(&mutex);
}</pre>
```

```
Algorithm 2 Active Barrier.
void activeBarrier()
{
    pthread_spin_lock(&lock);
    actNumArrived++;
    if(actNumArrived==NUM_THREADS) {
        actLoopFlag = false;
    }
    pthread_spin_unlock(&lock);
    while(actLoopFlag) {}
}
```

- Implemented two different barriers
 - "Passive" barrier
 - Mutexes and conditional wait statements
 - "Active" barrier
 - Spin locks and active polling



Barriers



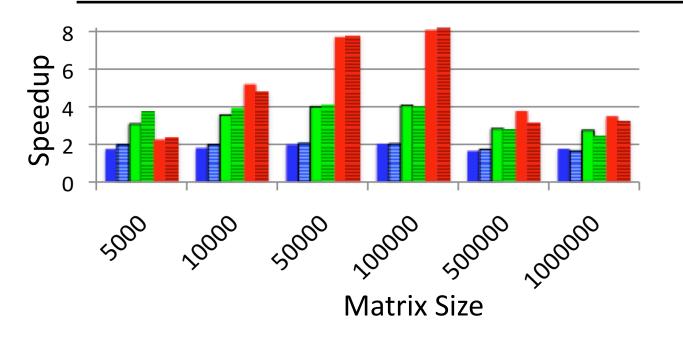
- Results for good data locality matrices
- Active/aggressive barriers essential for scalability



- Studied the importance of thread affinity
- Thread affinity allows threads to be pinned to cores
 - Less likely for threads to be switched (beneficial for cache utilization)
 - Ensures that threads are running on same socket



Thread Affinity

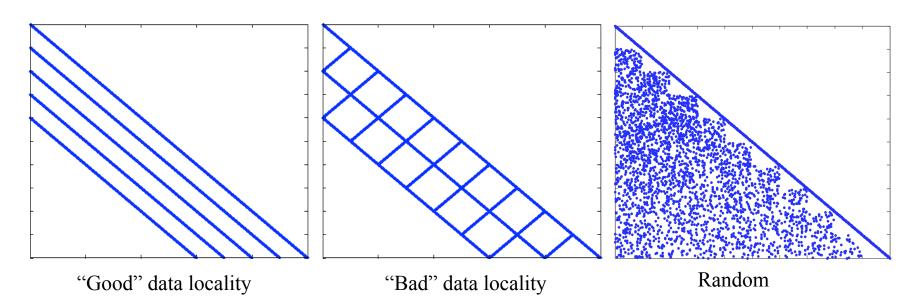


- NTA, threads=2
- TA, threads=2
- NTA, threads=4
- TA, threads=4
- NTA, threads=8
- TA, threads=8

- Results for good data locality matrices, active barrier
- Thread affinity not as important as active barrier
 - But can be beneficial for some problem sizes



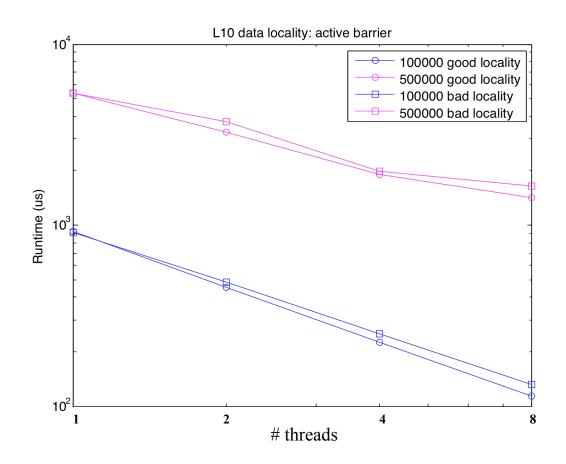
Factor 3: Data Locality



- Examined three different types of matrices
 - Same number of rows per level
 - Same number of nonzeros per row
- Allowed us to explore how data locality affects performance



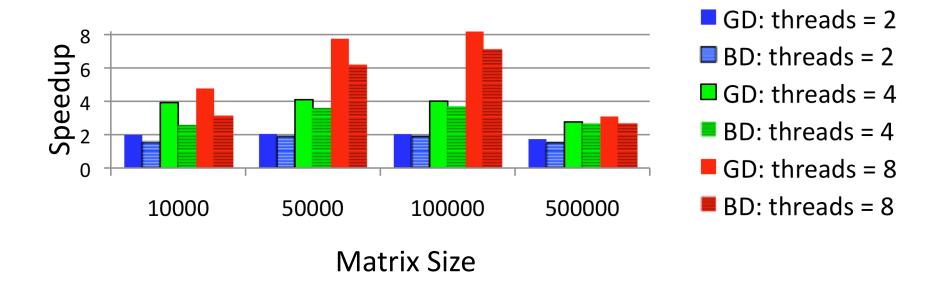
Data Locality: Good vs. Bad



- Results for good (GD) vs. bad data (BD) locality matrices
- Active barrier



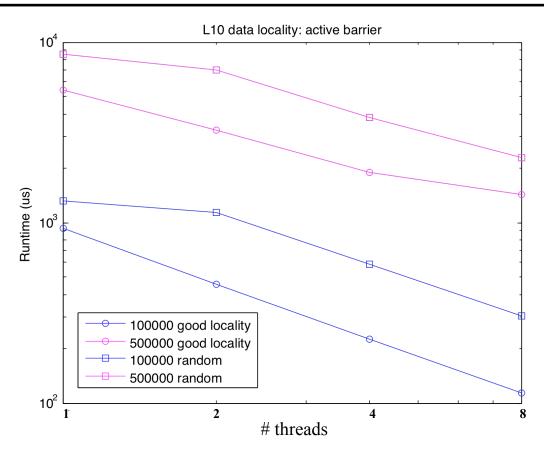
Data Locality: Good vs. Bad



- Results for good (GD) vs. bad data (BD) locality matrices
- Active Barrier



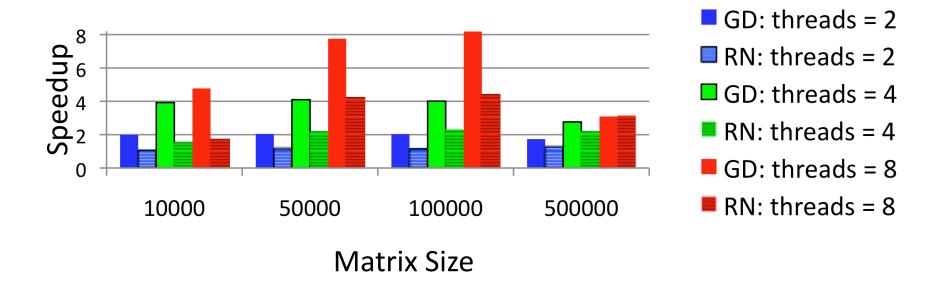
Data Locality: Good vs. Random



- Results for good data locality vs. random matrices
- Active barrier



Data Locality: Good vs. Random



- Results for good data locality (GD) vs. random (RN) matrices
- Active Barrier



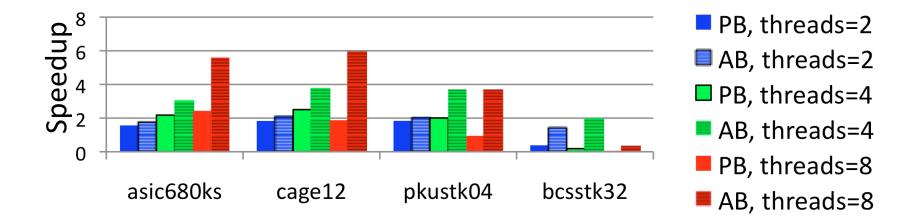
More Realistic Problems

Name	Ν	nnz	N / nlevels	Application area
asic680ks	682,712	2,329,176	13932.9	circuit simulation
cage12	130,228	2,032,536	1973.2	DNA electrophoresis
pkustk04	55,590	4,218,660	149.4	structural engineering
bcsstk32	44,609	2,014,701	15.1	structural engineering

- Symmetric matrices
- Incomplete Cholesky factorization (no fill)
- Average size of level important



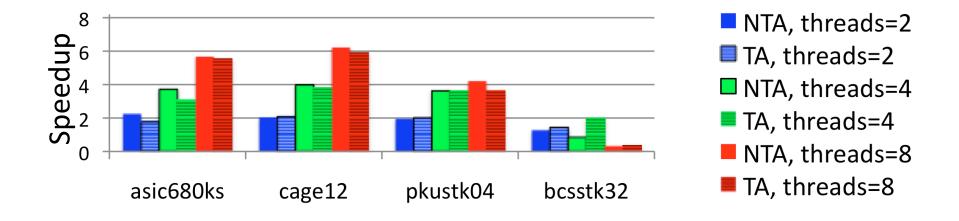
Realistic Problems: Barriers



- Problems with larger average level size scale fairly well
- Active/aggressive barrier important



Realistic Problems: Thread Affinity

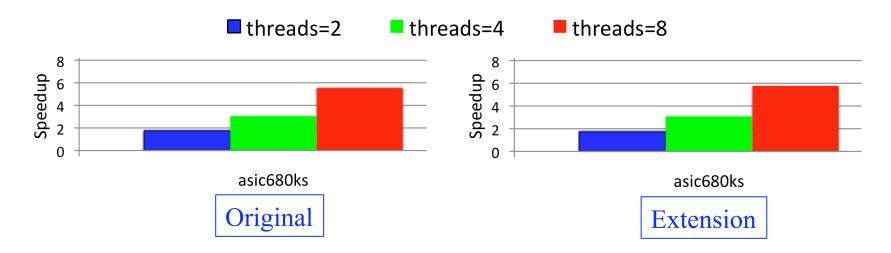


- Problems with larger average level size scale fairly well
- Thread affinity not particularly important



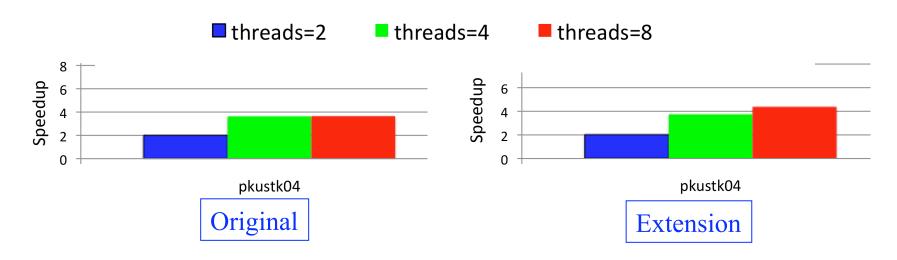
- Algorithm scales when average level size is high
- Couple factors hurt performance for small average level size
 - Many levels, many synchronization points
 - Not enough work in small levels (barrier cost significant)
- Implemented simple extension to address these problems
 - Serialize small levels below a certain threshold
 - Merge consecutive serialized levels
 - Reducing levels reduces synchronization points





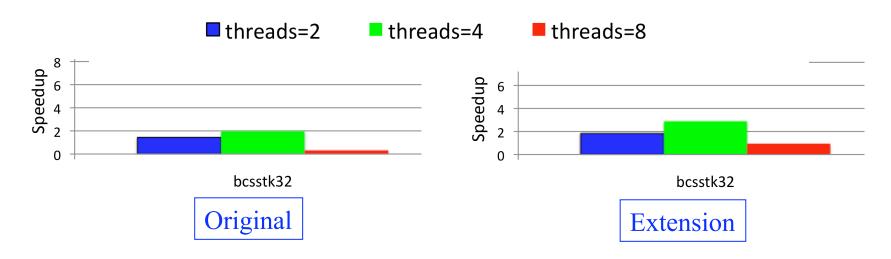
- Very slight improvement for problem that scale well
 - Not many small levels
 - Can reduce speedup if too aggressive in serialization





- Slight improvement for problem that originally did not scale quite so well
 - More small levels





- Significant improvement for problem that originally did not scale well
 - Many small levels
 - Great reduction in synchronization points
- Still does not scale well for 8 threads



- Presented threaded triangular solve algorithm
 - Level scheduling algorithm
- Studied impact of three factors on performance
 - Barrier type most important
- Good scalability for simple matrices and two realistic problems
- Scalability related to average level size
 - Simple extension to improve results when level sizes are small
 - Better algorithms needed for matrices with small average level size
- Algorithms being implemented in Trilinos
 - http://trilinos.sandia.gov