# Charm++ on the road to Exascale

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PEEPS @ VecPar
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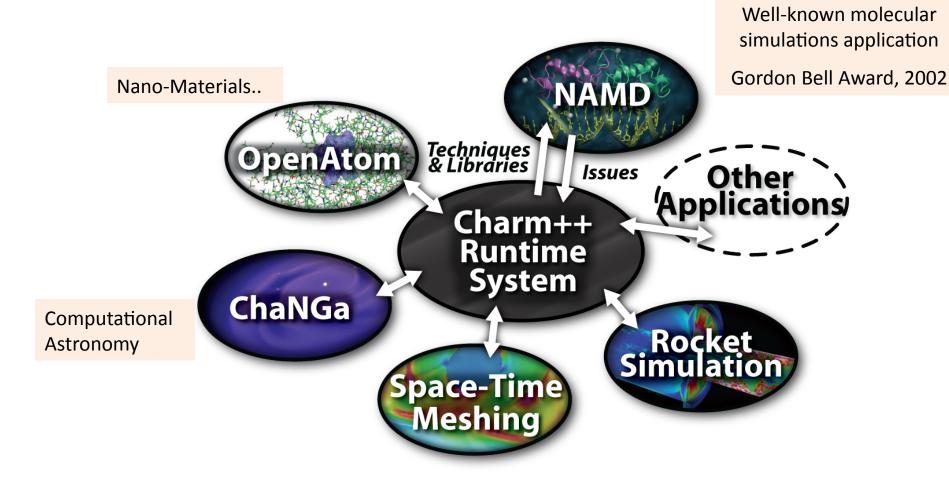




#### Overview

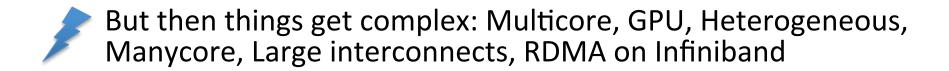
- Context
- History
- Philosophy
- Over-decomposition into Migratable Objects
- Some Current Research Efforts
  - Scalable Performance Analysis
  - Automatic Performance Tuning
  - New types of Runtime System Adaptation
  - Load Balancing
  - Topology Mapping
  - Accelerators
  - Incomplete but Useful Languages
- Promising Features in Charm++ for the Future

## Context (so far)



## A Glance at History

- 1987: Chare Kernel arose from parallel Prolog work
  - Dynamic load balancing for state-space search, Prolog, ...
- 1992: Charm++
- 1996: Charm++ in almost current form
  - Chare arrays, Measurement Based Dynamic Load balancing
- 1997: AMPI



- 2008-2011: Blue Waters:
  - Charm++ and NAMD will scale to 1PFlop/s sustained application performance

## PPL Mission and Approach

- To enhance <u>Performance and Productivity</u> in programming complex parallel applications
  - Performance: scalable to >100s of thousands of processors
  - Productivity: of human programmers
  - Complex: irregular structure, dynamic variations
- Approach: <u>Application Oriented yet CS centered</u> <u>research</u>
  - Develop enabling technology, for a wide collection of apps.
  - Develop, use and test it in the context of real applications

## Some Guiding Principles

- No magic
  - Parallelizing compilers have achieved close to technical perfection, but are not enough
  - Sequential programs obscure too much information
- Seek an optimal division of labor between the system and the programmer
- Design abstractions based solidly on use-cases
  - Application-oriented yet computer-science centered approach

L. V. Kale, "Application Oriented and Computer Science Centered HPCC Research", Developing a Computer Science Agenda for High-Performance Computing, New York, NY, USA, 1994, ACM Press, pp. 98-105.

### Migratable Objects (aka Processor Virtualization)

<u>Programmer:</u> [Over] decomposition into virtual processors / objects

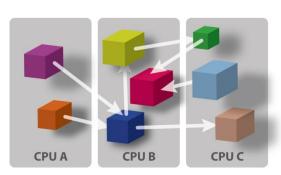
Runtime: Assigns VPs to processors

Enables adaptive runtime strategies

Implementations: Charm++, AMPI



**User View** 

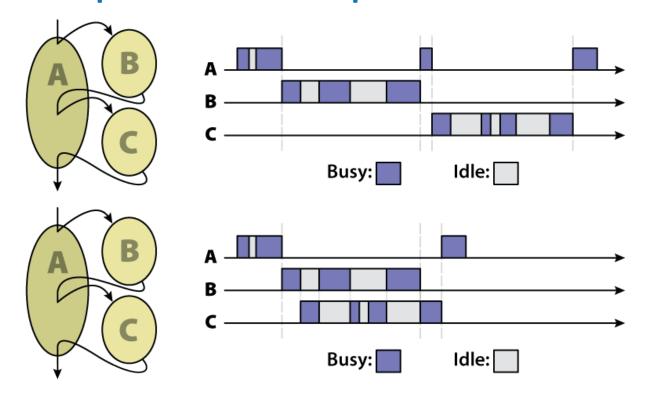


System View

#### **Benefits**

- Software engineering
  - Number of virtual processors can be independently controlled
  - Separate VPs for different modules
- Message driven execution
  - Adaptive overlap of communication
  - Predictability:
    - Automatic out-of-core
  - Asynchronous reductions
- Dynamic mapping
  - Heterogeneous clusters
    - Vacate, adjust to speed, share
  - Automatic checkpointing
  - Change set of processors used
  - Automatic dynamic load balancing
  - Communication optimization

## Adaptive overlap and modules



#### SPMD and Message-Driven Modules

(From A. Gursoy, Simplified expression of message-driven programs and quantification of their impact on performance, Ph.D Thesis, Apr 1994)

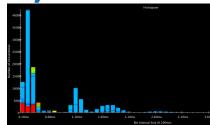
Modularity, Reuse, and Efficiency with Message-Driven Libraries: Proc. of the Seventh SIAM Conference on Parallel Processing for Scientific Computing, San Fransisco, 1995

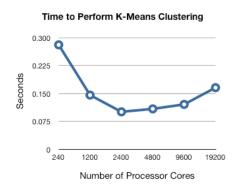
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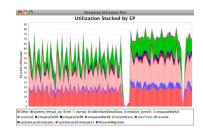
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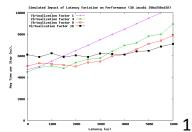
Scalable Performance Analysis

- Programming models must support scalable performance analysis or automatic tuning!
- Scalable performance analysis idioms
  - Must do in parallel
  - Use end-of-run when machine is available to you
  - E.g. parallel k-means clustering
- Live streaming of performance data
  - stream live performance data out-of-band in user-space to enable powerful analysis idioms



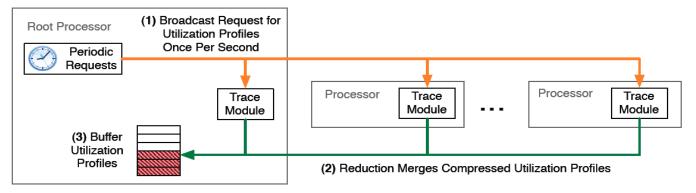




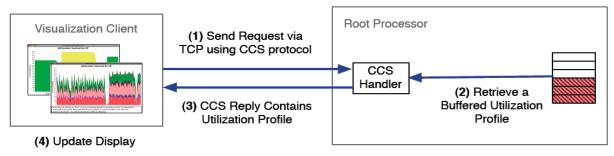


## Live Streaming System Overview

- Interleave/Compose performance monitoring with application execution
  - A) Gathering Performance Data in Parallel Runtime System:

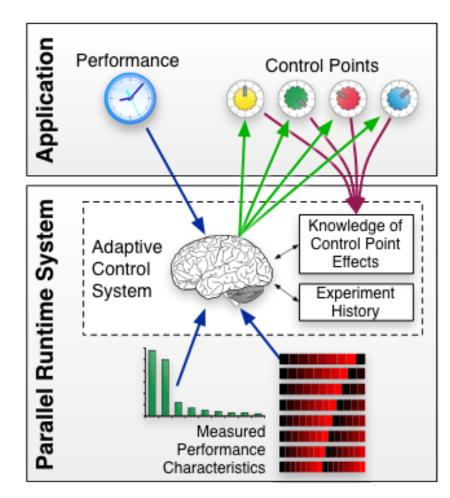


B) Visualizing Performance Data:



## **Automatic Performance Tuning**

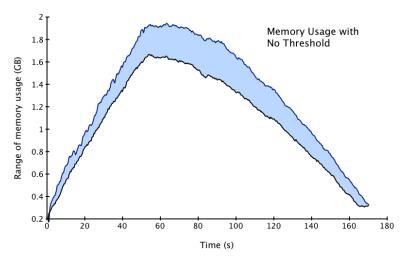
- The runtime system dynamically reconfigures applications
- Tuning/Steering is based on runtime observations :
  - Idle time, overhead time, grain size, # messages, critical paths, etc.
- Applications expose in structured manner: tunable parameters AND information about the parameters

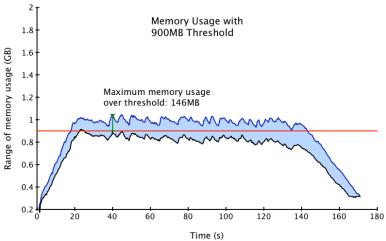


Isaac Dooley, and Laxmikant V. Kale, *Detecting and Using Critical Paths at Runtime in Message Driven Parallel Programs*, 12th Workshop on Advances in Parallel and Distributed Computing Models (APDCM 2010) at IPDPS 2010.

#### New Types of Dynamic Adaptation

- Memory Aware Scheduling
- The Charm++ scheduler was modified to adapt its behavior.
- It can give preferential treatment to annotated entry methods when available memory is low.
- The memory usage for an LU Factorization program is reduced, enabling further scalability.





Isaac Dooley, Chao Mei, Jonathan Lifflander, and Laxmikant V. Kale, A Study of Memory-Aware Scheduling in Message Driven Parallel Programs, PPL Technical Report 2010

## Scalability & Performance Summary

- We think dynamic adaptation is good.
- Automatic Adaptation Is Necessary
  - Manual tuning of even simple applications requires super-experts!
  - There are dynamically changing performance characteristics of application and of the system.
  - Need good support for intelligent automatic tuning?
  - Need good support for partially automated performance analysis?

#### Need Composability of:

- Applications
- Performance Measuring
- Intelligent Autotuning
- Runtime system adaptation

- ...

## Large Scale Parallelism: Load Balancing at Petascale

- Our older load balancing strategies don't scale on extremely large machines
  - Consider an application with 1M objects on 64K processors

#### Centralized

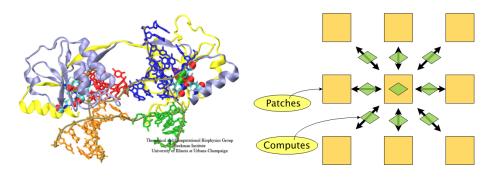
- Object load data are sent to processor 0
- Integrate to a complete object graph
- Migration decision is broadcast from processor 0
- Global barrier

#### Distributed

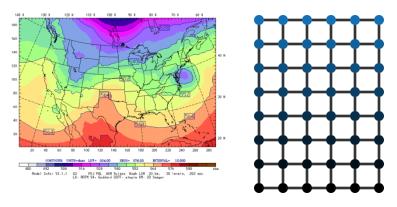
- Load balancing among neighboring processors
- Build partial object graph
- Migration decision is sent to its neighbors
- No global barrier
- Topology-aware
  - On 3D Torus/Mesh topologies

## Mapping Objects onto Machine Interconnect Topology

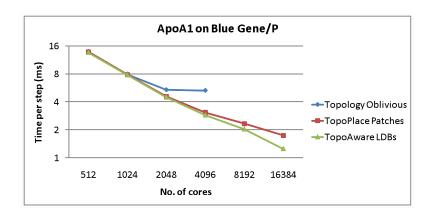
#### **Charm++ Applications**



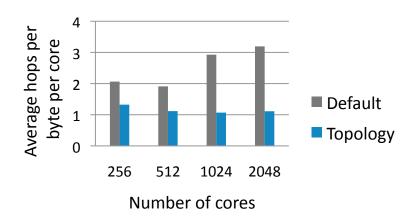
**MPI** Applications



Molecular Dynamics - NAMD



Weather Research & Forecasting Model



## **Load Balancing Summary**

- Need appropriate load balancing algorithms
- Many concerns:
  - Need excellent quality load balance (comm. & compute)
  - Machine interconnect topology
  - Complex hierarchy of processor cores
  - Cost of performing each load balancing operation
- Charm++ model, due to its encouraging of overdecomposing, provides many load balancing opportunities.

## Accelerators and Heterogeneity

- GPUs, Larrabee, IBM Cell processor, ..
- It turns out that some of the Charm++ features are a good fit for these
  - Charm++ model already decomposes applications into small somewhat self-contained pieces.
- With minor changes to Charm++, tighter encapsulations of objects helps bridge barriers to multiple-memory-space platforms (or non-cache coherent systems)

Kunzman and Kale, Towards a Framework for Abstracting Accelerators in Parallel Applications: Experience with Cell, finalist for best student paper at SC09

## Incomplete But Useful Languages

- Why use just one(or two) languages/paradigms?
- High level parallel languages:
  - Charisma: static data flow
  - SDAG: static & dynamic data flow
  - Multiphase Shared Arrays: shared arrays with access modes
  - Charj: Java-like version of Charm++ that supports better compiler generation of utility code (efficient object serialization, ...)
- Multiple paradigms coexist with each other and can be interleaved in time (message driven scheduling)!

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## Promising Features in Charm++ for the Future

#### 1. Overdecomposition

- Portability to exotic architectures, better encapsulation
- Greater ability to automatically manage (load balancing, fault tolerance, ...)

#### 2. Efficient composability of multiple modules

 application, paradigms, libraries, performance analysis, autotuning, debugging, ...

#### 3. Automatic management where possible

Adaptation within runtime system, and in application

### Exascale?

- Exascale machines will look different.
  - Power consumption, exotic/new architectures, ...
- Charm++ looks like a promising candidate with sufficient margins & flexibility to fit onto a wide range of conceivable architectures.
- Charm++ model is complex, but then again so are many other choices (which require super-experts to achieve good performance):
  - PGAS+threads+cuda+...
  - MPI+OpenMP+OpenCL+...

## **Final Summary**

- Emerging parallel machines:
  - Big
  - Differing architectural designs
  - Increasingly hard to program
- Charm++ model has some interesting and promising features that address some of these difficulties:
  - Overdecomposition
  - Efficient composition

## Questions?

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